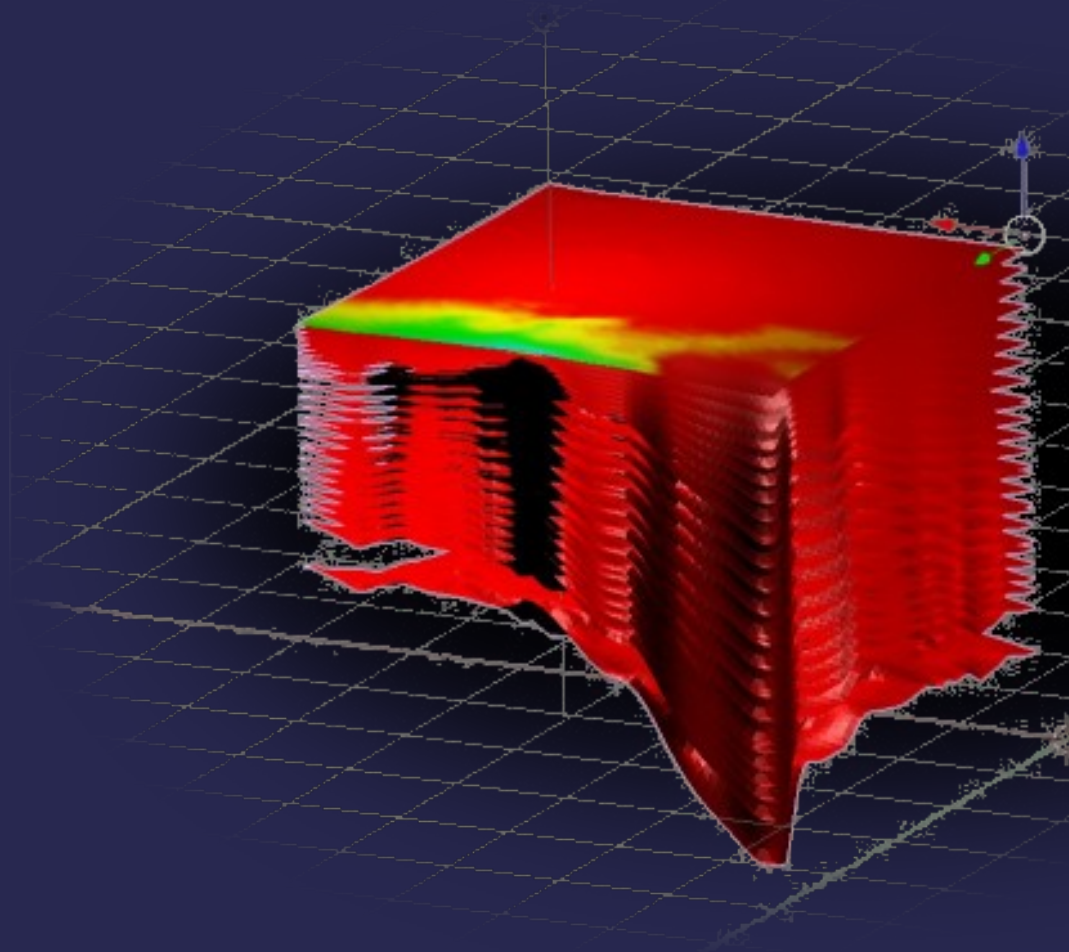


# Parallel Computations and the Finite Element Method

Joachim B Haga

Sintef

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The equilibrium

$$\underline{\nabla} \cdot \underline{\underline{\sigma}} + \underline{\underline{g}} = 0,$$

$$\underline{\underline{\sigma}} = \lambda \underline{\underline{\nabla}} \cdot \underline{\underline{u}} \underline{\underline{I}} + \mu (\underline{\underline{\nabla}} \underline{\underline{u}} + \underline{\underline{\nabla}} \underline{\underline{u}}^T) - \beta_s (3\lambda + 2\mu)$$

and  $(\lambda, \mu)$  as the Lamé constants

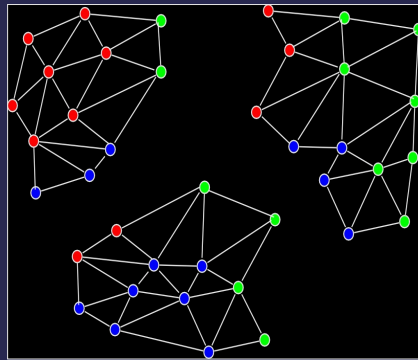
Following Darcy's law

$$c \frac{\partial T}{\partial t} + \underline{\underline{g}}_T + \underline{\underline{c}}_T \underline{\underline{\nabla}}_D \cdot \underline{\underline{v}}_D = 0$$

$$\underline{\underline{v}}_D = \phi \underline{\underline{v}}_s = -\frac{k}{\mu} \underline{\underline{\nabla}} T$$



## Overview of parallel computing, and why we should care

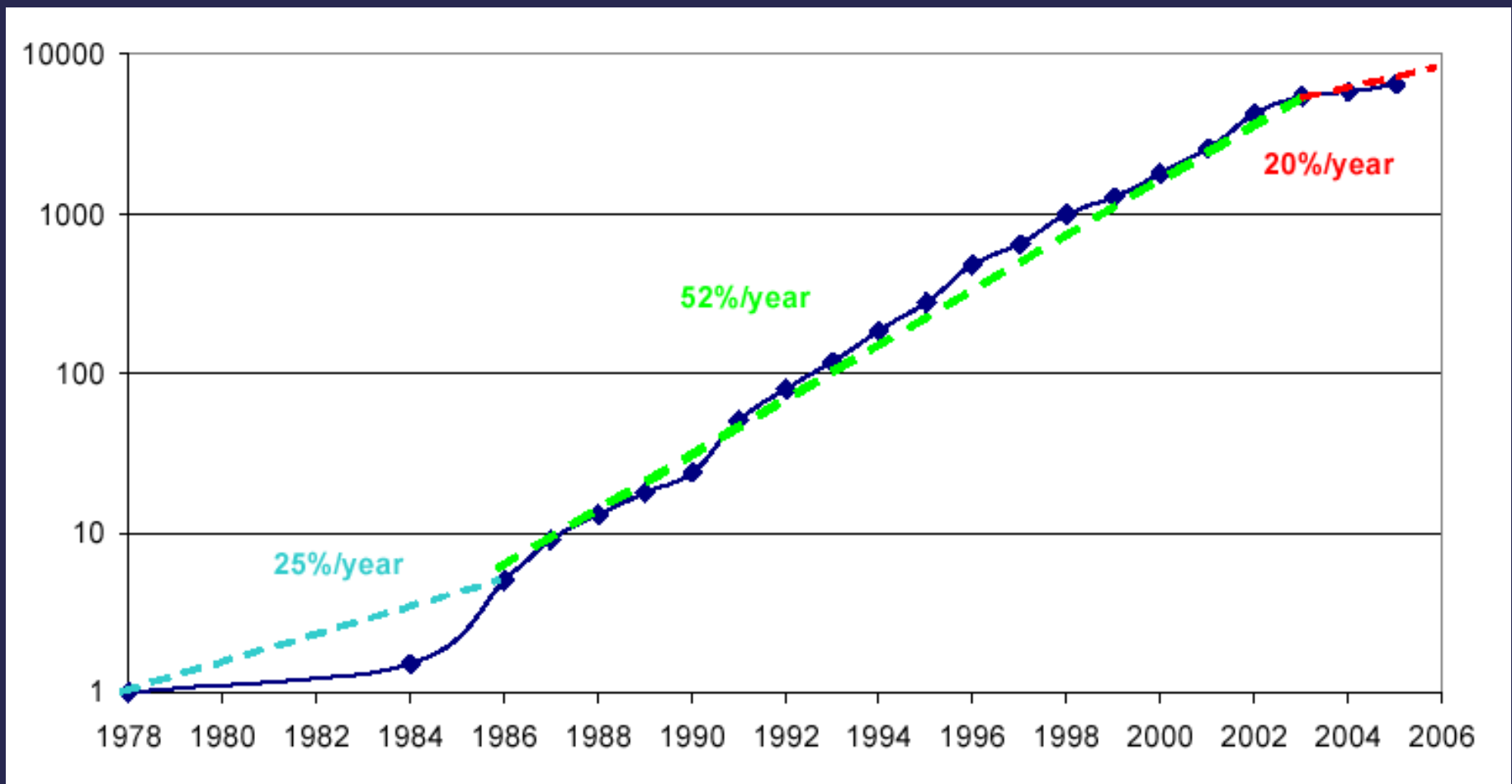


## Parallelisation of the finite element method

$$\begin{aligned} A^0 &\rightarrow P^1 \\ A^1 &= R^1 A^0 P^1 \rightarrow P^2 \\ &\vdots \\ A^e &= R^e A^{e-1} P^e \rightarrow P^{e+1} \\ &\vdots \\ A^L &= R^L A^{L-1} P^L \end{aligned}$$

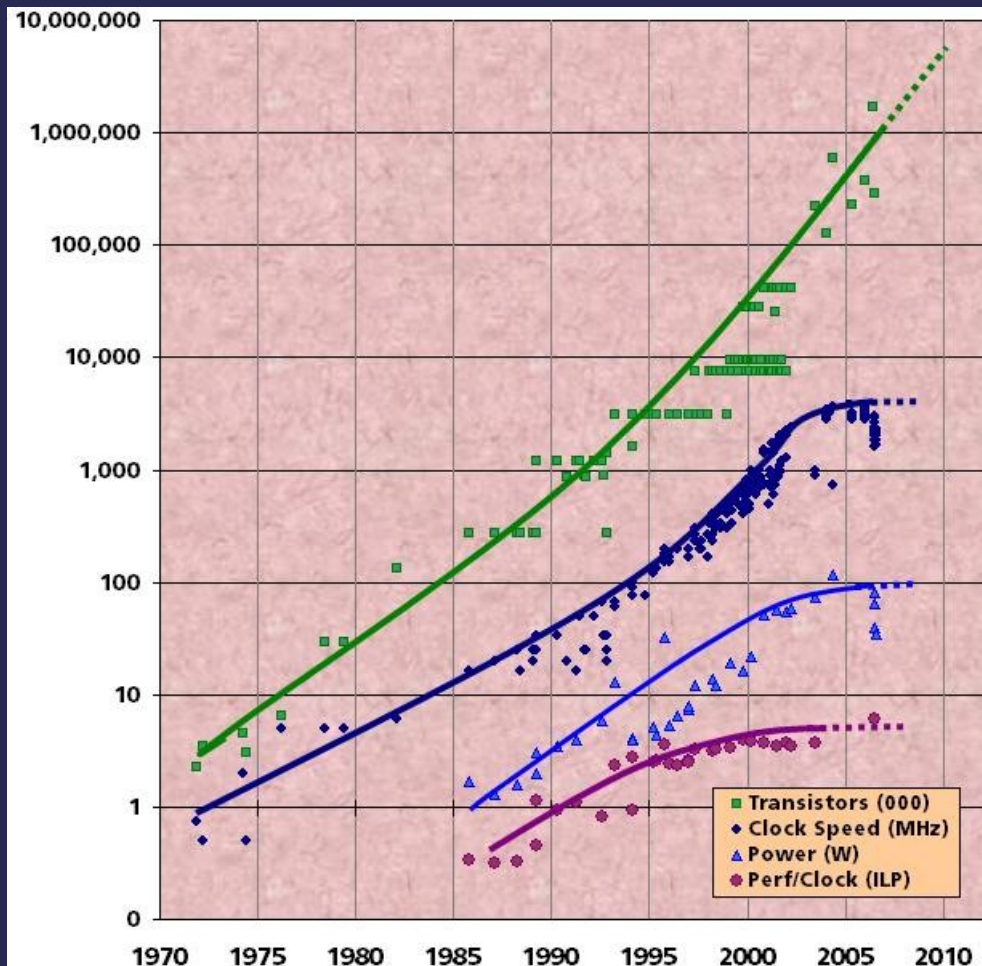
## Algebraic multigrid as a parallel preconditioner

# Single-processor performance has stalled since ~2002 ...



(Hennessy and Patterson, 2006)

... so the future is parallel



(Olokotun and Sutter)

- Power wall
- ILP wall
- Memory wall
- VLSI wall



Lower clock speeds  
Simpler modules  
Multiple cores

# A number of alternatives exist for parallel computations

Distributed memory (**MPI**)

Shared memory (**OpenMP**)

GPU programming (CUDA, CTM)

Cell (PS3)

SIMD vector engines (Cray etc.)

- There is no consensus on the paradigm for the parallel future
  - so stick with established standards

# OpenMP vs MPI example: inner product

```
double inner(vec a, vec b)
{
    double sum = 0;

    for (int i=0; i<a.size(); i++)
        sum += a[i]*b[i];

    return sum;
}
```

# OpenMP vs MPI example: OpenMP version (transparent to caller)

```
double inner(vec a, vec b)
{
    double sum = 0;

    # pragma omp parallel for \
        private(i) \
        reduction(+:sum)
    for (int i=0; i<a.size(); i++)
        sum += a[i]*b[i];

    return sum;
}
```

## OpenMP vs MPI example: MPI version (requires caller knowledge)

```
double inner(vec a, vec b)
{
    double sum = 0, glob_sum;

    for (int i=0; i<n_local_rows; i++)
        sum += a[i]*b[i];

    MPI_Allreduce(&sum, &glob_sum, 1, MPI_DOUBLE,
                 MPI_SUM, MPI_COMM_WORLD);
    return glob_sum;
}
```



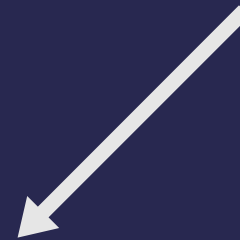
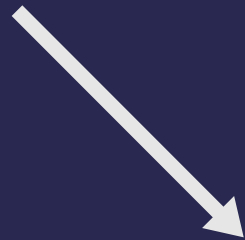
# OpenMP, MPI, hybrid: pros and cons

## OpenMP

- shared memory
- allows gradual parallelisation
- does little for the “memory wall”

## MPI

- distributed memory
- allows clusters
- not transparent



## Hybrid

- OpenMP + MPI
- scales better than either alone

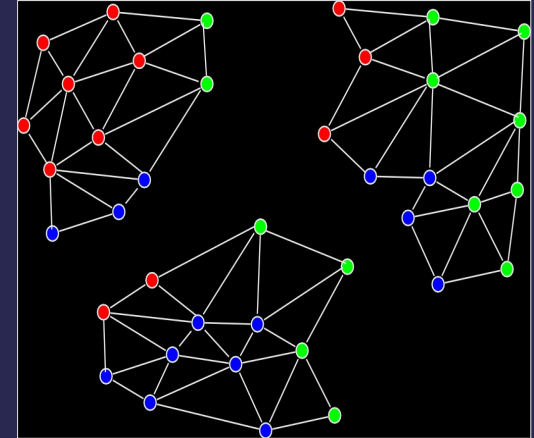
# Parallelising the Finite Element Method

→ Grid partitioning

→ Parallel assembly

→ Parallel linear algebra

→ Parallel I/O

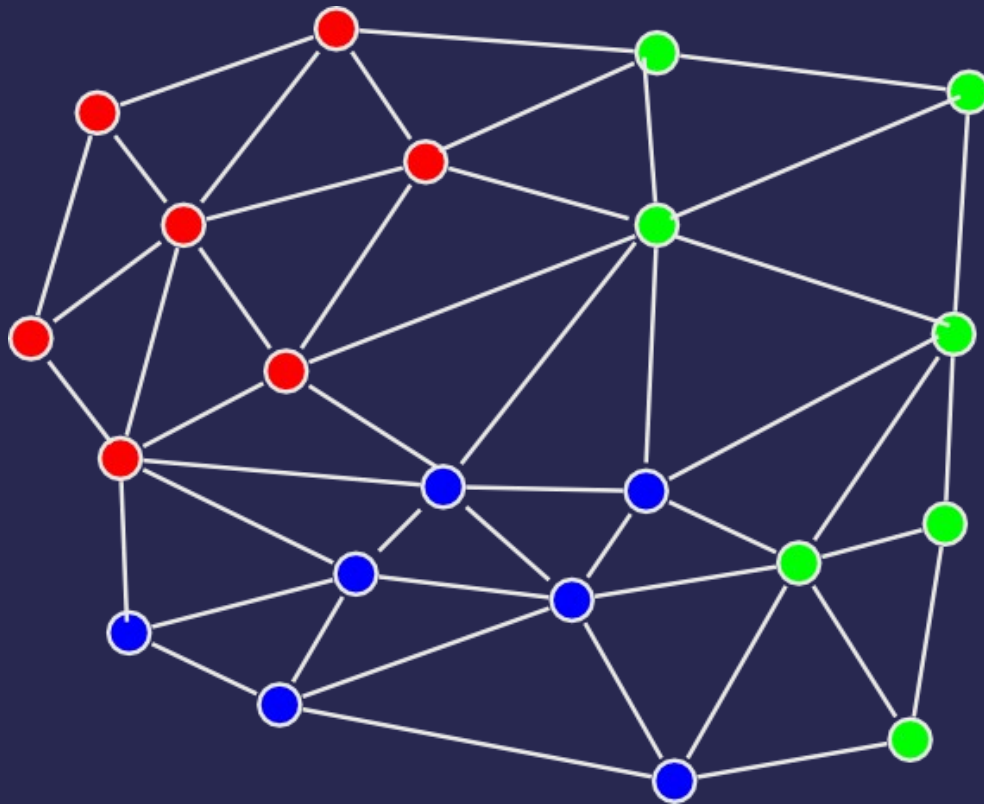


# Nodal grid partitioning



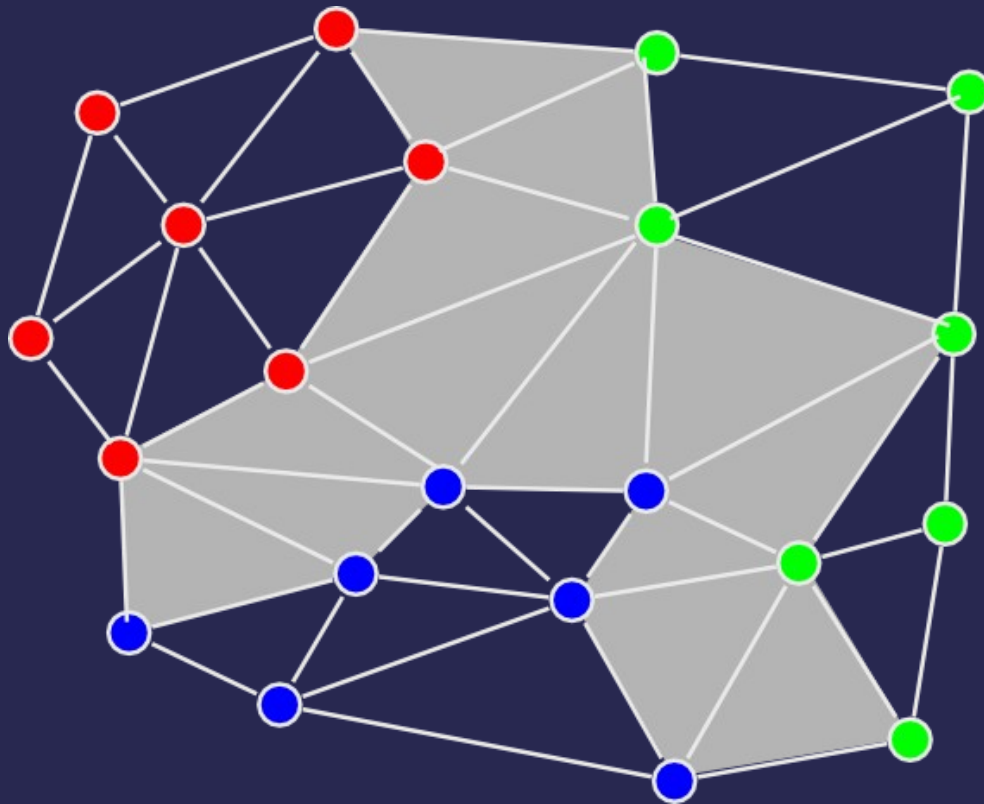
→ Construct a graph from the mesh

# Nodal grid partitioning



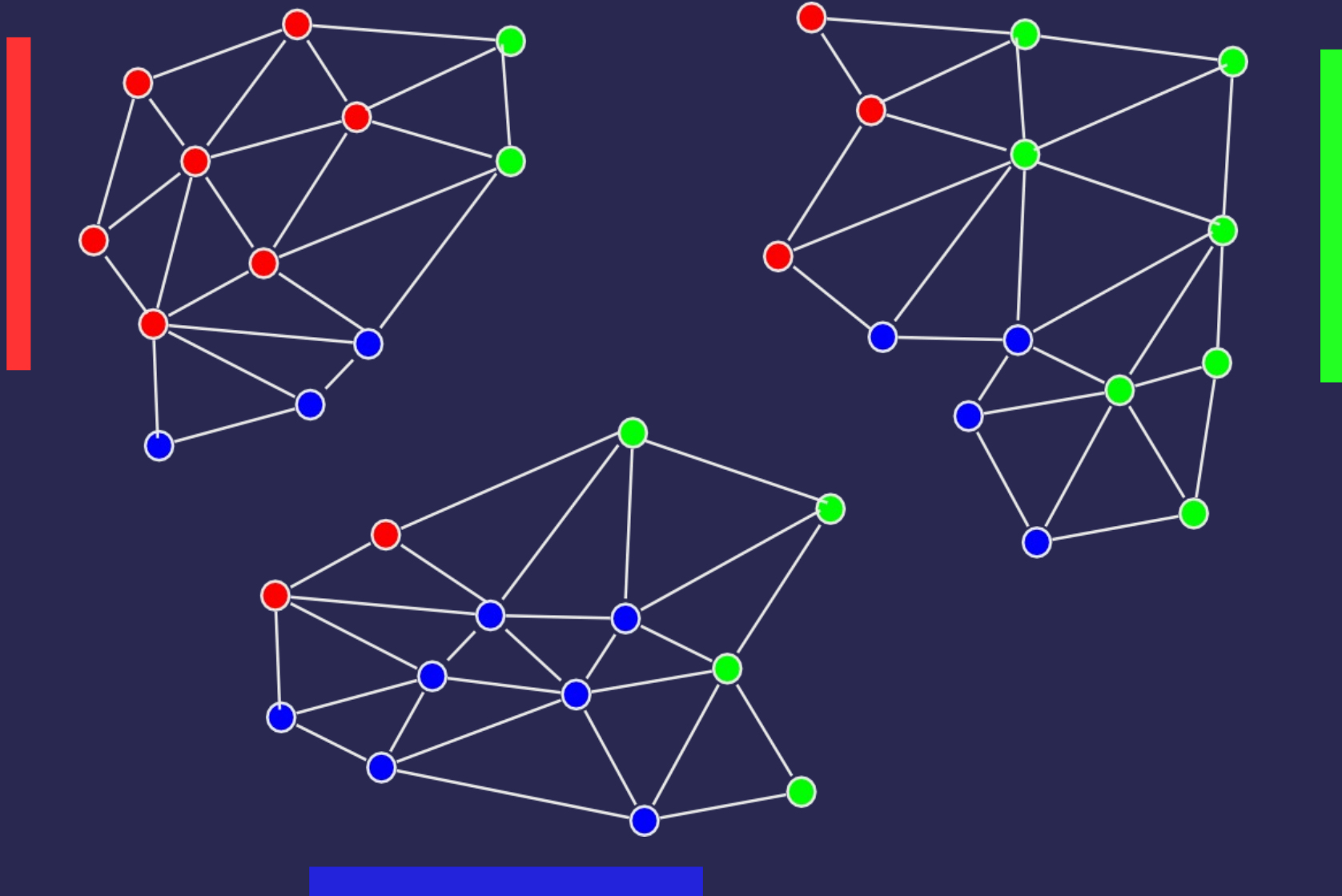
- Construct a graph from the mesh
- Partition the graph
  - METIS, ParMETIS
  - Scotch, PScotch

# Nodal grid partitioning



- Construct a graph from the mesh
- Partition the graph
  - METIS, ParMETIS
  - Scotch, PScotch
- An element is on a cpu if any of its nodes are
- Result: A shared band of border elements

# Nodal grid partitioning



The advantage of nodal grid partitioning is that each matrix row is complete on one processor

Dual grid partitioning (of the elements) creates a shared band of *nodes* instead of *elements*

- fewer nodes overall
- slightly smaller communication cost
- the shared nodes have no canonical placement

But with nodal partitioning, every node is the responsibility of exactly one processor

- the matrix row associated with a node is in one place
- makes e.g. algebraic multigrid much easier

# The major components of a finite element solver

## Parallel assembly

- trivial, just assemble locally

## Parallel linear algebra

- requires care
- preconditioning remains a problem

## Parallel I/O

- commonly to local disk
- gather result post-process



# Distributed memory parallel linear algebra

## Vector addition

- no communication

## Vector inner products, norms

- exclude ghost-nodes from local norm
- reduction (sum) operation after local norm

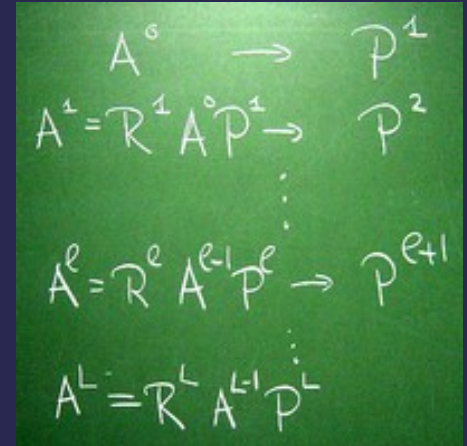
## Matrix-vector product

- update ghost-node values before multiplication
- (optionally) update again after multiplication

## Preconditioning

- hard...

# Parallel preconditioning


$$\begin{aligned} A^0 &\rightarrow P^1 \\ A^1 &= R^1 A^0 P^1 \rightarrow P^2 \\ &\vdots \\ A^l &= R^l A^{l-1} P^l \rightarrow P^{l+1} \\ &\vdots \\ A^L &= R^L A^{L-1} P^L \end{aligned}$$

- A few preconditioners are easy to parallelise (Jacobi, for example, is trivial)
- But many popular ones are not (e.g. ILU)
- Geometric multigrid is possible, but that is not an option on unstructured grids
- So what about algebraic multigrid?

# Algebraic multigrid phases

## Coarsening

- problem with coarsening across processor boundaries
- can operate decoupled (sub-optimal)
- various coupling strategies
- limits coarse grid size (to #cpus)

## Projection / interpolation

- requires no communication

## Smoothing

- Jacobi smoothing is decoupled
- Gauss-Seidel most popular?

# Changes needed to support parallel operations for Trilinos/ML are minor (with nodal partitioning)

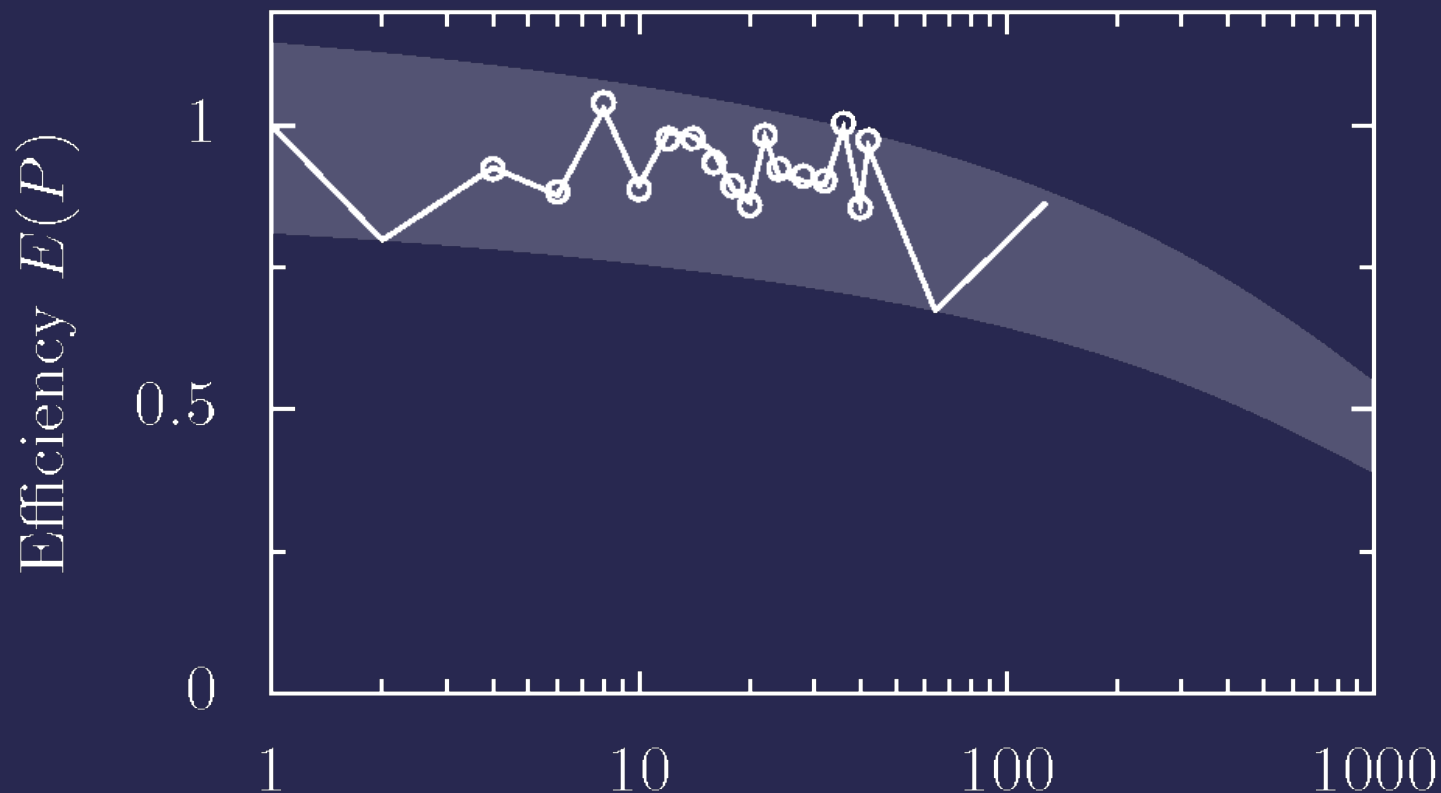
In addition to the sequential interface code, we need to pass in

- the number of local rows
- a function which updates ghost-nodes in a vector

Also, the “matvec” and “apply” functions must of course be parallel-aware

- “matvec” updates ghost-nodes before multiplication
- “apply” updates ghost-nodes after V cycle

# Algebraic multigrid shows great promise as parallel preconditioner for the coupled system



Parallel efficiency of the AMG-preconditioned BiCGStab solver (ML/Trilinos+Diffpack)

# Some key questions

## What is the target?

- |                       |          |
|-----------------------|----------|
| Multicore, < 8 cpus   | → OpenMP |
| Cluster, or > 16 cpus | → MPI    |
| > 100 cpus            | → Hybrid |
| > 1000 cpus           | → ???    |

## Can existing libraries be used?

- Trilinos (MPI)
- Hypre (hybrid)



# The geomechanical model: Equation for the fluid pressure

$$S \frac{\partial p}{\partial t} = \nabla \cdot (\Lambda \nabla p) - \nabla \cdot (\Lambda \rho_f (1 - \beta_f (T - T_0)) \mathbf{g})$$

$S$  storage coefficient

$\Lambda$  mobility of flow

$\rho$  density

$\beta$  thermal expansion coefficient

# The geomechanical model: Equation for the temperature

$$C \frac{\partial T}{\partial t} + \rho_f C_f \mathbf{v}_D \cdot \nabla T = \nabla \cdot (\kappa \nabla T)$$

$\mathbf{v}_D$  Darcy velocity (in porous media)  
 $= \phi \mathbf{v}_f = -\Lambda(\nabla p - \rho_f(1 - \beta_f(T - T_0))\mathbf{g})$

$\phi$  porosity

$C$  bulk heat capacity  
 $= \phi \rho_f C_f + (1 - \phi) \rho_s C_s$

$C_f$  specific heat

$\kappa$  thermal conductivity



# The geomechanical model: Equation for the deformation (elastic)

$$0 = \nabla \cdot \sigma + \rho \mathbf{g}$$

$$\sigma = (\lambda \nabla \cdot \mathbf{u} - \alpha p - \beta_s(3\lambda + 2\mu)(T - T_0)) \mathbf{I} + 2\mu \epsilon$$

$\mu, \lambda$  Lamé material constants

$\alpha$  Biot factor  
 $\approx 1$

$\mathbf{u}$  displacement field

$\epsilon$  deformation tensor  
 $= (\nabla \mathbf{u} + (\nabla \mathbf{u})^T) / 2$